Arithmetic Flags and Instructions

Chapter 6

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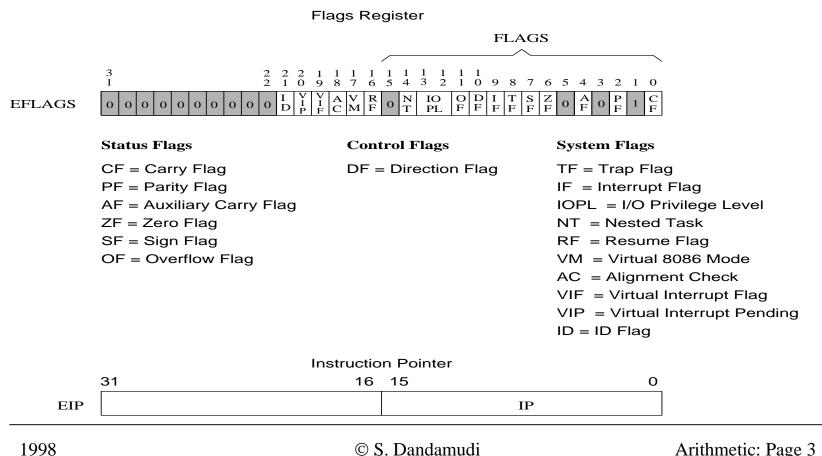
Outline

- Status flags
 - * Zero flag
 - * Carry flag
 - * Overflow flag
 - * Sign flag
 - * Auxiliary flag
 - * Parity flag
- Arithmetic instructions
 - * Addition instructions
 - * Subtraction instructions
 - * Multiplication instructions
 - * Division instructions

- Application examples
 - * PutInt8
 - * GetInt8
- Multiword arithmetic
 - * Addition
 - * Subtraction
 - * Multiplication
 - * Division
- Performance: Multiword multiplication

Status Flags

Six status flags monitor the outcome of arithmetic, logical, and related operations



- Status flags are updated to indicate certain properties of the result
 - * Example: If the result is zero, zero flag is set
- Once a flag is set, it remains in that state until another instruction that affects the flags is executed
- Not all instructions affect all status flags
 - * add and sub affect all six flags
 - * inc and dec affect all but the carry flag
 - * mov, push, and pop do not affect any flags

Example

```
; initially, assume ZF = 0
              ; ZF is still zero
       AL,55H
mov
       AL,55H; result is 0
sub
                ; ZF is set (ZF = 1)
                ; ZF remains 1
push
       BX
                ; ZF remains 1
       BX,AX
mov
                  ZF remains 1
       DX
pop
       CX,0
                ; ZF remains 1
mov
inc
       CX
                ; result is 1
                  ZF is cleared (ZF = 0)
```

Zero Flag

- * Indicates zero result
 - If the result is zero, ZF = 1
 - Otherwise, ZF = 0
- * Zero can result in several ways (e.g. overflow)

- » All three examples result in zero result and set ZF
- * Related instructions

jz jump if zero (jump if ZF = 1)

jnz jump if not zero (jump if ZF = 0)

- Uses of zero flag
 - * Two main uses of zero flag
 - » Testing equality
 - Often used with cmp instruction

```
cmp char, '\$'; ZF = 1 if char is \$
```

cmp AX, BX

- » Counting to a preset value
 - Initialize a register with the count value
 - Decrement it using dec instruction
 - Use jz/jnz to transfer control

Consider the following code

```
sum := 0
for (I = 1 to M)
  for (j = 1 to N)
    sum := sum + 1
  end for
end for
```

Assembly code

```
sub
        AX,AX ; AX := 0
         DX,M
  mov
outer_loop:
         CX,N
  mov
inner_loop:
  inc
         \mathbf{A}\mathbf{X}
  loop
         inner_loop
  dec
         DX
  jnz
         outer loop
exit_loops:
         sum, AX
  mov
```

- Two observations
 - * loop instruction is equivalent to

```
dec CX
jnz inner loop
```

- » This two instruction sequence is more efficient than the **loop** instruction (takes less time to execute)
- » loop instruction does not affect any flags!
- * This two instruction sequence is better than initializing CX to zero and executing

```
inc CX
cmp CX,N
jle inner_loop
```

Carry Flag

- * Records the fact that the result of an arithmetic operation on *unsigned* numbers is out of range
- * The carry flag is set in the following examples

```
mov AL,0FH mov AX,12AEH add AL,0F1H sub AX,12AFH
```

* Range of 8-, 16-, and 32-bit unsigned numbers

size	range
8 bits	0 to $255(2^8-1)$
16 bits	0 to 65,535 $(2^{16}-1)$
32 bits	0 to $4,294,967,295$ ($2^{32}-1$)

- * Carry flag is not set by inc and dec instructions
 - » The carry flag is *not set* in the following examples

mov	AL,OFFH	mov	AX,0
inc	AL	dec	AX

* Related instructions

jc jump if carry (jump if CF = 1)

jnc jump if no carry (jump if CF = 0)

* Carry flag can be manipulated directly using

stc set carry flag (set CF to 1)

clc clear carry flag (clears CF to 0)

cmc complement carry flag (inverts CF value)

- Uses of carry flag
 - * To propagate carry/borrow in multiword addition/subtraction

 $1 \leftarrow \text{carry from lower } 32 \text{ bits}$

x = 3710 26A8 1257 9AE7H

y = 489B A321 FE60 4213H

7FAB C9CA 10B7 DCFAH

- * To detect overflow/underflow condition
 - » In the last example, carry out of leftmost bit indicates overflow
- * To test a bit using the shift/rotate instructions
 - » Bit shifted/rotated out is captured in the carry flag
 - » We can use jc/jnc to test whether this bit is 1 or 0

Overflow flag

- * Indicates out-of-range result on *signed* numbers
 - Signed number counterpart of the carry flag
- * The following code sets the overflow flag but not the carry flag

```
mov AL,72H ; 72H = 114D
add AL,0EH ; 0EH = 14D
```

* Range of 8-, 16-, and 32-bit signed numbers

size	range	
8 bits	- 128 to +127	2^7 to $(2^7 - 1)$
16 bits	- 32,768 to +32,767	2^{15} to $(2^{15} - 1)$
32 bits	-2,147,483,648 to $+2,147,483,647$	2^{31} to $(2^{31} - 1)$

- Signed or unsigned: How does the system know?
 - * The processor does not know the interpretation
 - * It sets carry and overflow under each interpretation

Unsigned interpretation

mov AL,72H

add AL, OEH

jc overflow

no overflow:

(no overflow code here)

. . . .

overflow:

(overflow code here)

. . . .

Signed interpretation

mov AL,72H

add AL, OEH

jo overflow

no_overflow:

(no overflow code here)

. . . .

overflow:

(overflow code here)

. . . .

* Related instructions

jo jump if overflow (jump if OF = 1)

jno jump if no overflow (jump if OF = 0)

* There is a special software interrupt instruction

into interrupt on overflow

Details on this instruction in Chapter 12

- Uses of overflow flag
 - * Main use
 - » To detect out-of-range result on signed numbers

Sign flag

- * Indicates the sign of the result
 - Useful only when dealing with signed numbers
 - Simply a copy of the most significant bit of the result

* Examples

mov	AL,15	mov	AL,15
add	AL,97	sub	AL,97
clears	the sign flag as	sets th	ne sign flag as
the resi	ult is 112	the res	sult is –82
(or 0111000 in binary)		(or 10	101110 in binary)

* Related instructions

js jump if sign (jump if SF = 1) jns jump if no sign (jump if SF = 0)

- Usage of sign flag
 - * To test the sign of the result
 - * Also useful to efficiently implement countdown loops
- Consider the count down loop: The count down loop can be implemented as **for** (i = M downto 0)<loop body> CX,M mov end for for_loop: If we don't use the jns, we <loop body> need cmp as shown below: dec CX cmp jns for_loop for_loop jl

Auxiliary flag

- * Indicates whether an operation produced a carry or borrow in the low-order 4 bits (nibble) of 8-, 16-, or 32-bit operands (i.e. operand size doesn't matter)
- * Example

» As there is a carry from the lower nibble, auxiliary flag is set

* Related instructions

- » No conditional jump instructions with this flag
- » Arithmetic operations on BCD numbers use this flag

aaa ASCII adjust for addition

aas ASCII adjust for subtraction

aam ASCII adjust for multiplication

aad ASCII adjust for division

daa Decimal adjust for addition

das Decimal adjust for subtraction

Chapter 11 has more details on these instructions

* Usage

» Main use is in performing arithmetic operations on BCD numbers

Parity flag

- * Indicates even parity of the low 8 bits of the result
 - PF is set if the lower 8 bits contain even number 1 bits
 - For 16- and 32-bit values, only the least significant 8 bits are considered for computing parity value

* Example

```
mov AL,53 53D = 0011 0101B add AL,89 89D = 0101 1001B 142D = 1000 1110B
```

» As the result has even number of 1 bits, parity flag is set

* Related instructions

```
jump on even parity (jump if PF = 1)

jump on odd parity (jump if PF = 0)
```

- * Usage of parity flag
 - » Useful in writing data encoding programs
 - » Example: Encodes the byte in AL (MSB is the parity bit)

```
parity_encode PROC
    shl AL
    jp parity_zero
    stc
    jmp move_parity_bit
    parity_zero:
        clc
    move_parity_bit:
        rcr AL
parity_encode ENDP
```

Arithmetic Instructions

- Pentium provides several arithmetic instructions that operate on 8-, 16- and 32-bit operands
 - » Addition: add, adc, inc
 - » Subtraction: sub, sbb, dec, neg, cmp
 - » Multiplication: mul, imul
 - » Division: div, idiv
 - » Related instructions: cbw, cwd, cdq, cwde, movsx, movzx
 - * There are few other instructions such as **aaa**, **aas**, etc. that operate on decimal numbers
 - » See Chapter 11 for details

Addition instructions

* Basic format

add destination, source

» Performs simple integer addition

destination := destination + source

* Five operand combinations are possible

add register, register

add register, immediate

add memory, immediate

add register, memory

add memory, register

* Basic format

adc destination, source

» Performs integer addition with carry

destination := destination + source + CF

- * Useful in performing addition of long word arithmetic
- The three carry flag manipulating instructions are useful

stc set carry flag (set CF to 1)

clc clear carry flag (clears CF to 0)

cmc complement carry flag (inverts CF value)

- * The final instruction **inc** requires a single operand inc destination
 - » Performs increment operation

destination := destination + 1

- » The operand is treated as an unsigned number
- * Does not affect the carry flag
 - » Other five status flags are updated
- * In general

inc BX

is better than

add BX,1

» Both take same time but **inc** version takes less space

Subtraction instructions

sub destination, source

» Performs simple integer subtraction

destination := destination - source

sbb destination, source

» Performs integer subtraction with borrow

destination := destination - source - CF

dec destination

» Performs decrement operation

destination := destination - 1

• Subtraction instructions (cont'd)

neg destination

» Performs sign reversal

destination := 0 - destination

» Useful in signed number manipulation

cmp destination, source

- » Performs subtraction without updating destination destination source
- » Updates all six status flags to record the attributes of the result
- » The cmp instruction is typically followed by a conditional jump instruction

Multiplication

- * More complicated than add/sub
 - » Produces double-length results
 - E.g. Multiplying two 8 bit numbers produces a result that requires 16 bits
 - » Cannot use a single multiply instruction for signed and unsigned numbers
 - add and sub instructions work both on signed and unsigned numbers
 - For multiplication, we need separate instructions

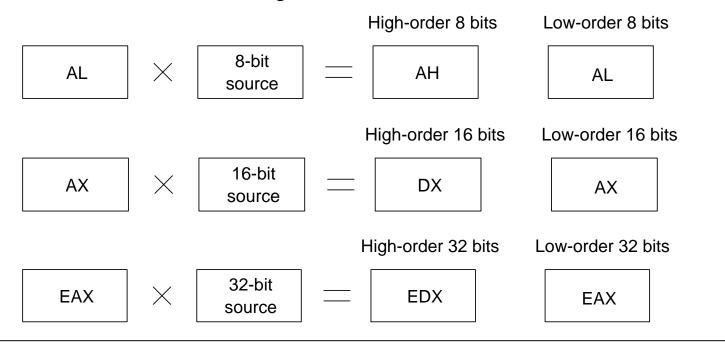
mul for unsigned numbers

imul for signed numbers

Unsigned multiplication

mul source

» Depending on the **source** operand size, the location of the other source operand and destination are selected



* Example

```
mov AL,10
mov DL,25
mul DL
```

produces 250D in AX register (result fits in AL)

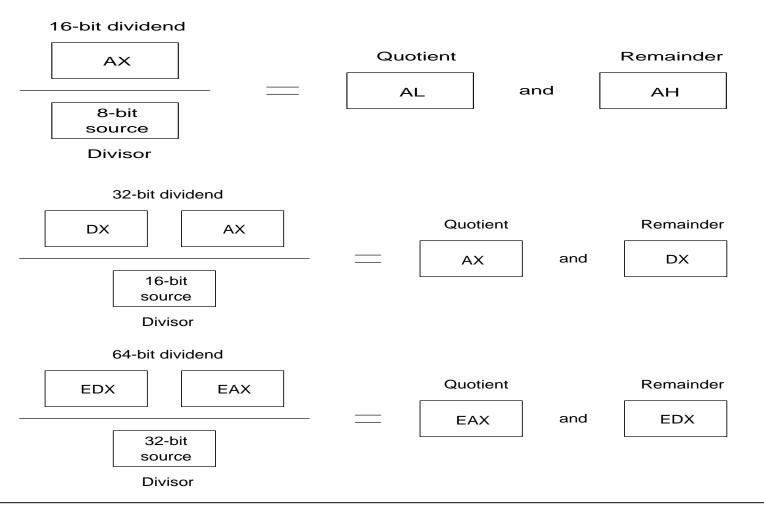
- The **imul** instruction can use the same syntax
 - » Also supports other formats
 - * Example

```
mov DL,0FFH ; DL := -1
mov AL,0BEH ; AL := -66
mul DL
```

produces 66D in AX register (again, result fits in AL)

- Division instruction
 - * Even more complicated than multiplication
 - » Produces two results
 - Quotient
 - Remainder
 - » In multiplication, using a double-length register, there will not be any overflow
 - In division, divide overflow is possible
 - → Pentium provides a special software interrupt when a divide overflow occurs
 - * Two instructions as in multiplication

div	source	for unsigned numbers
idiv	source	for signed numbers



Example

```
mov AX,00FBH ; AX := 251D
mov CL,0CH ; CL := 12D
div CL
```

produces 20D in AL and 11D as remainder in AH

Example

produces 17D in AX and 47D as remainder in DX

- Signed division requires some help
 - » We extended an unsigned 16 bit number to 32 bits by placing zeros in the upper 16 bits
 - » This will not work for signed numbers
 - To extend signed numbers, you have to copy the sign bit into those upper bit positions
 - * Pentium provides three instructions in aiding sign extension
 - » All three take no operands

cbw converts byte to word (extends AL into AH)

cwd converts word to doubleword (extends AX into DX)

cdq converts doubleword to quadword

(extends EAX into EDX)

- * Some additional related instructions
 - » Sign extension

cwde converts word to doubleword (extends AX into EAX)

» Two move instructions

movsx dest,src (move sign-extended src to dest)
movzx dest,src (move zero-extended src to dest)

- » For both move instructions, **dest** has to be a register
- » The **src** operand can be in a register or memory
 - If src is 8-bits, dest has to be either a 16 bit or 32 bit register
 - If **src** is 16-bits, **dest** has to be a 32 bit register

Example

```
mov AL,0A1H ; AL := -95D
cbw ; AH = FFH
mov CL,0CH ; CL := 12D
idiv CL
```

produces –7D in AL and –11D as remainder in AH

Example

```
mov AX,0EBE5 ; AX := -5147D cwd ; DX := FFFFH mov CX,012CH ; CX := 300D idiv CX
```

produces –17D in AX and –47D as remainder in DX

Application Examples

• PutInt8 procedure

* To display a number, repeatedly divide it by 10 and display the remainders obtained

	quotient	remainder
108/10	10	8
10/10	1	0
1/10	0	1

- * To display digits, they must be converted to their character form
 - » This means simply adding the ASCII code for zero (see line 24)

line 24: add AH,'0'

Application Examples (cont'd)

• **GetInt8** procedure

- * To read a number, read each digit character
 - » Convert to its numeric equivalent
 - » Multiply the running total by 10 and add this digit

Input digit	Numeric	Number := $Number*10 + N$
	value (N)	
Initial value		0
'1'	1	0*10+1=1
'5'	5	1*10+5=15
' 8'	8	15 * 10 + 8 = 158

Multiword Arithmetic

- Arithmetic operations (add, sub, mul, and div) work on 8-, 16-, or 32-bit operands
- Arithmetic on larger operands require multiword arithmetic software routines
- Addition/subtraction
 - * These two operations are straightforward to extend to larger operand sizes
 - * Need to use adc/sbb versions to include the carry generated by the previous group of bits
 - * Example addition on the next slide

```
; Adds two 64-bit numbers in EBX: EAX and EDX: ECX.
:The result is returned in EBX:EAX.
:Overflow/underflow conditions are indicated
; by setting the carry flag.
;Other registers are not disturbed.
 add64 PROC
       add
               EAX, ECX
       adc
               EBX, EDX
       ret
  add64
```

Multiplication

- * We consider two algorithms
 - » Longhand multiplication
 - Uses the method that we are familiar with
 - Needs addition operations only
 - Examines each bit in the multiplier and adds the multiplicand if the multiplier bit is 1
 - → Appropriate shifting is required
 - » Using the mul instruction
 - Chops the operand into 32-bit chunks and applies mulinstruction
 - Similar to the addition example seen before

* Longhand multiplication

Final 128-bit result in P:A P := 0; count := 64 A := multiplier; B := multiplicand while (count > 0) if (LSB of A = 1) then P := P+B CF := carry generated by P+B else CF := 0 end if

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count := count-1

end while

shift right CF:P:A by one bit position

- * Using the **mul** instruction
 - » A 64-bit number is treated as two 32-bit numbers
 - A is considered as consisting of A1A0 (similarly B)
 - Left shift operation replaces zeros on the right

```
temp := A0 X B1
temp := left shift temp
    by 32 bits
result := result + temp
temp := A1 X B1
temp := left shift temp
    by 32 bits
result := result + temp
```

Division

- * To implement n-bit division (A by B), we need an additional n+1 bit register P
- * Core part of the algorithm
 - » Test the sign of P
 - » if P is negative
 - left shift P:A by one bit position
 - P := P + B
 - » else
 - left shift P:A by one bit position
 - -P := P B

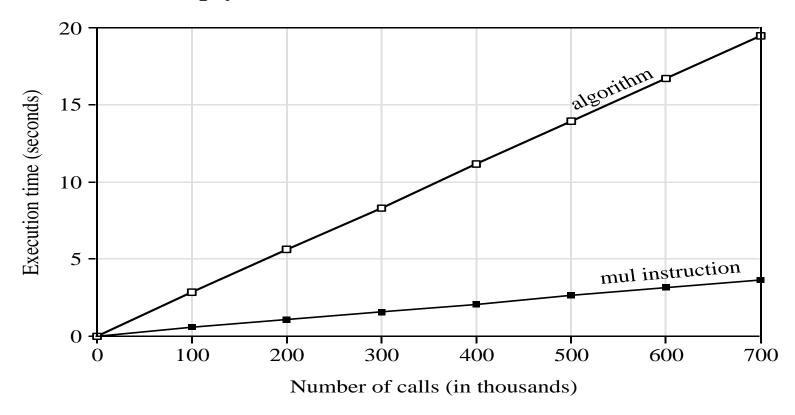
Division Algorithm

```
P := 0; count := 64
A := dividend
B := divisor
while (count > 0)
  if (P is negative)
  then shift left P:A
       by 1 bit position
       P := P + B
  else shift left P:A
       by 1 bit position
       P := P-B
  end if
```

```
A = quotient, P = remainder
  <u>if</u> (P is negative)
  then set low-order
          bit of A to 0
  else set low-order
          bit of A to 1
  end if
  count := count-1
end while
<u>if</u> (P is negative)
  P := P+B
end if
```

Performance: Multiword Multiplication

- Longhand version versus mul version
 - * To multiply $2^{64}-1 \times 2^{64}-1$



Performance: Multiword Multiply (cont'd)

- Using add versus mul instruction
 - * Certain special cases of multiplication can be done by a series additions (e.g. power of 2 by shift operations)
 - * Example: Multiplication by 10
 - * Instead of using **mul** instruction, we can multiply by 10 using **add** instructions as follows (performs $AL \times 10$):

```
sub AH,AH ; AH := 0
mov BX,AX ; BX := x
add AX,AX ; AX := 2x
add AX,AX ; AX := 4x
add AX,BX ; AX := 5x
add AX,AX ; AX := 10x
```

Performance: Multiword Multiply (cont'd)

• Multiplication of 2³²–1 by 10

